

# **Predicting Web Service Levels During VM Live Migrations**



Helmut Hlavacs and Thomas Treutner

Research Group Entertainment Computing, University of Vienna, Austria

#### Introduction

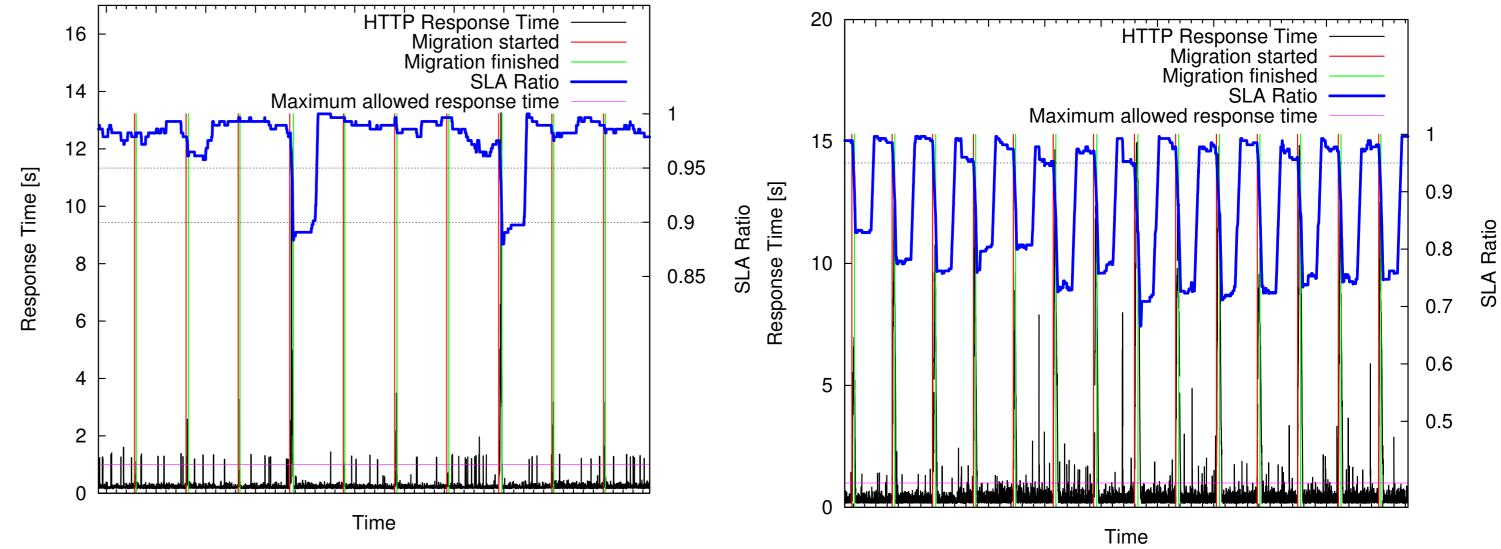
- VM live migration important for energy efficiency
- Enables us to establish energy efficient target distribution of VMs Supposedly no perceivable service downtime while live migrating
- Live migration is resource intensive (iterative page copying)
- **Experiments:** Influence on service levels while migrating?
- Modelling: Predict service levels based on utilization?

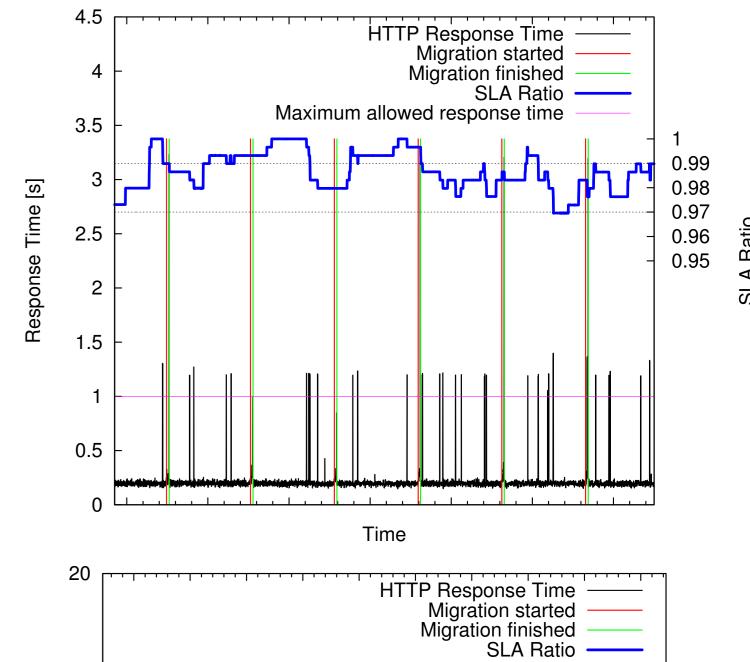
## **Scenario**

- Virtualized data center, static consolidation (P2V)
- Provisioning for peak load, still bad energy efficiency
- ► E.g., 9-5-cycles, very low utilization at night

#### Service Levels for Low/Medium/High Workload Scenarios

- **Low** (top right): Slightly increased response times during live migration, seldom response time violations
- ► Medium (bottom left): SLA ratio generally satisfies the 97% limit
- **High** (bottom right): Often and heavy violations, unacceptable low service levels, typically decreased by 20-25% percentage points



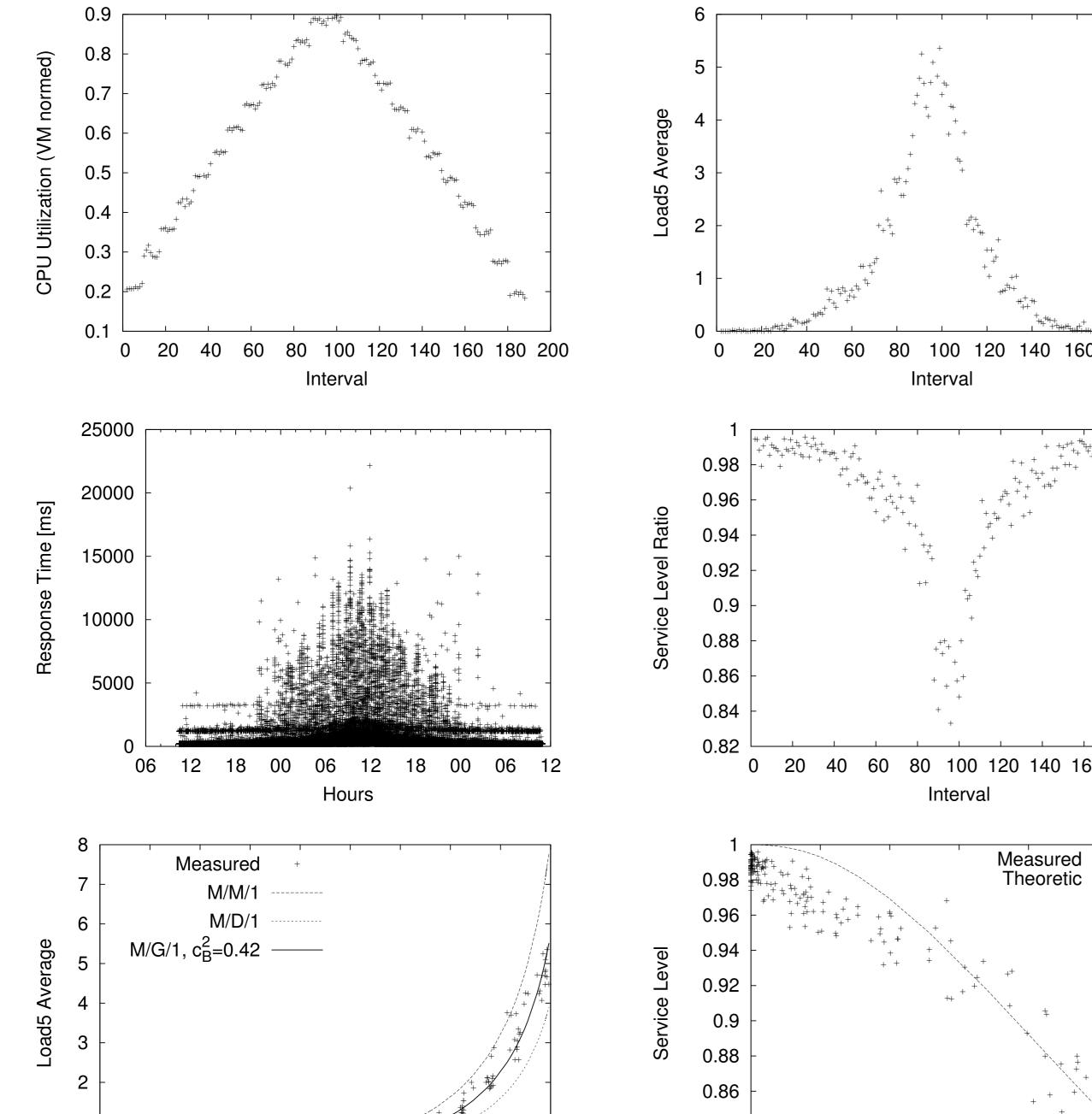


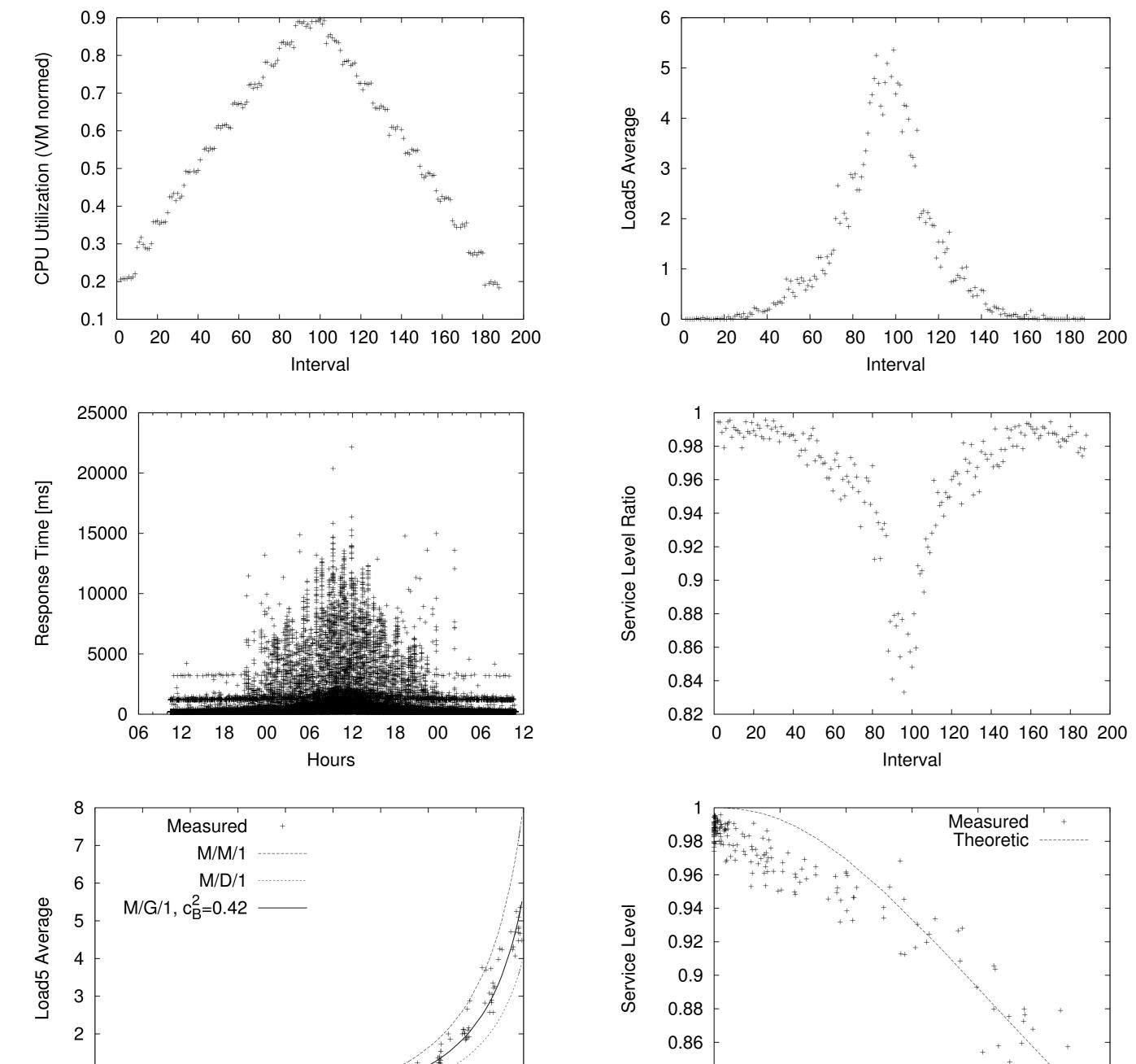
- **Live migration enables dynamic consolidation**
- **But: Seldom used, fear of possible side effects**
- $\blacktriangleright \Rightarrow$  Identify and quantify effects on (web) service levels  $\blacktriangleright \Rightarrow$  Find most influencing utilization metrics

#### Experiment

- **Two servers, a single VM, migrating forth and back**
- VM disk image on central node (Gbit, open-iscsi)
- qemu-kvm VM: Linux, Apache2, PHP5, MediaWiki
- SQL VM and load generation on an extra nodes
- **Logging utilization of servers and VM, more than 100 variables**
- ► Rise load from 50 to 600 concurrent virtual users and back
- ► Migrate every 15min, track response time of last 5min
- ► Maximum allowed response time: 1s

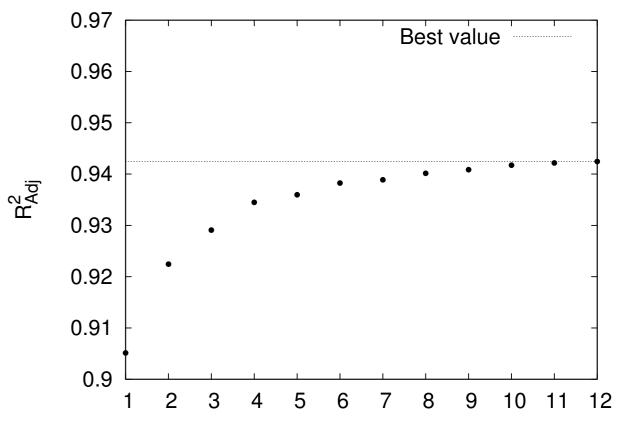
## **Data Overview**





# **Model Selection**

- Stepwise model selection: **Akaike Information Criterion**
- ► Finds trade-off between number of parameters (model size) and goodness of fit (model quality)
- **For comparison: Exhaustive** all-subsets-regression (LEAPS)
- ► LEAPS: Find best of all possible models for given range of model size
- Computationally intensive even if number of variables is limited
- $ightarrow 
  m R^2_{Adi}$  increases slightly with increased model complexity. UNIX load5 contributes  $\mathsf{R}^2_{\mathsf{Adi}}$  of  $\sim 90\%$



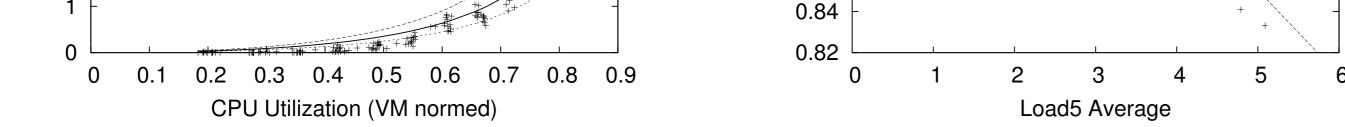
#### Model Size

#### **Most Influencing Model Parameters**

Variable	Meaning	Estimate	Std. Error	<b>Pr(</b> >  t )
Intercept		2.395e+00	5.069e-01	3.00e-06
wp01_load5	VM UNIX load5	-1.871e-02	2.627e-03	3.67e-12
wp01_swapUsed	VM swap used	-7.656e-07	8.809e-08	< 2e-16
wp01_residentSize_SQ	Suared amount of resident memory used by	-4.652e-14	1.652e-14	0.00506
	the qemu-kvm process.			
<pre>src_host_cpu_proc_s</pre>	Tasks created/s, source host.	-2.475e+00	9.166e-01	0.00716
<pre>src_host_cpu_proc_s_SQ</pre>		1.091e+00	4.133e-01	0.00856
wp01_cpu_util_vmnorm_SQ		-1.328e-01	3.187e-02	3.63e-05
wp01_cpu_util_vmnorm	CPU util measured inside VM	9.517e-02	2.316e-02	4.64e-05
wp01_load5_SQ	Squared UNIX load5 of the VM.	-1.140e-03	1.918e-04	5.22e-09
$wp01\_freeMemRatio\_SQ$	Squared ratio of free memory inside the VM.	1.976e-02	9.462e-03	0.03727

#### Conclusions

- Impact of live migration on SL depends on amount of workload
- **Tighter SLAs can be fulfilled during low and medium workload**
- Migrating during high load causes massive decrease of service level
- Service level variance during a live migration to 90% predictable using only a single variable, the UNIX load5 average, models with 12 variables can explain 95% of variance



- ► We can interpret the UNIX load as approximation to **Q**, the average number of jobs in a Markovian M/M/1 queue, and the VM's CPU utilization as ho, the system utilization:  $\overline{\mathbf{Q}}_{M/M/1} = \frac{\rho^2}{1-\rho}$
- UNIX load is exponentially averaged by definition and the service times are not necessarily exactly exponentially distributed:  $\overline{\mathbf{Q}}_{M/G/1} = \frac{\rho^2}{1-\rho} \cdot \frac{(1+c_B^2)}{2}$ For deterministic service times  $c_B^2 = 0$ , resulting in  $\overline{Q}_{M/D/1} = \frac{\rho^2}{2(1-\rho)}$ Simple linear regression delivers the coefficient  $c_B^2 = 0.42$  $\mathbf{P}(\mathbf{T} \leq \mathbf{x}) = \mathbf{F}_{\mathbf{T}}(\mathbf{x}) = \mathbf{1} - \mathbf{e}^{-\mu(1-\rho)\mathbf{x}}$ , the theoretical probability that the response time **T** is lower than or equal to a limit **x** for a given service rate  $\mu$
- Systems using live migration as a mechanism to realize a more energy efficient target distribution and have service level targets need to consider the UNIX load average, but typical hypervisors do not collect/export this information Hypervisors should be extended to export load information (cf. free memory)

#### **Future Work**

- Influence of additional VMs (idle, utilized, mixed)
- Linux and qemu-kvm: Kernel Samepage Merging (KSM)
- Database VM migration, currently taboo due to potentially severe influence
- qemu-kvm parameters: Bandwidth limits, maximum allowed downtime
- Predict migration delay, energy consumption, service downtime

Created with Latexbeamerposter http://www-i6.informatik.rwth-aachen.de/~dreuw/latexbeamerposter.php

<firstname>.<lastname>@univie.ac.at

http://cs.univie.ac.at/research/research-groups/entertainment-computing/